

MA 511: Computer Programming

Lecture 25

http://www.iitg.ernet.in/psm/indexing_ma511/y08/index.html

Partha Sarathi Mandal

psm@iitg.ernet.ac.in

Dept. of Mathematics, IIT Guwahati

Semester 1, 2008-09

Mon 10:00-10:55 Tue 11:00-11:55 Fri 9:00-9:55 Class: 1G2

MA512 Lab : Wed 14:00-16:55

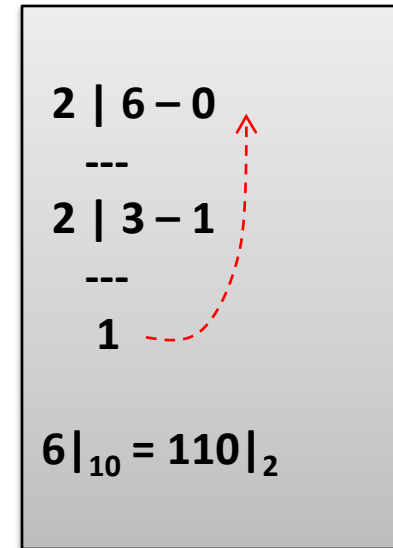
LOW-LEVEL PROGRAMMING

Number systems

<u>Decimal</u>	<u>Binary</u>	<u>Octal</u>	<u>Hexadecimal</u>
0	0000	0	0
1	0001	1	1
2	0010	2	2
3	0011	3	3
4	0100	4	4
5	0101	5	5
6	0110	6	6
7	0111	7	7
8	1000	10	8
9	1001	11	9
10	1010	12	A
11	1011	13	B
12	1100	14	C
13	1101	15	D
14	1110	16	E
15	1111	17	F

Decimal & Binary conversion

Decimal	Binary	Decimal
0	= 000	= $0 \times 2^2 + 0 \times 2^1 + 0 \times 2^0$
1	= 001	
2	= 010	
3	= 011	= $0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$
4	= 100	
5	= 101	
6	= 110	= $1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0$
7	= 111	= $1 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$



One's complement Operator (\sim)

- It's a unary operator that causes the bits of its operand to be inverted so that 1s become 0s and 0s become 1s.
- Operator always precedes its operand.
- The operand must be an integer-type quantity (including, long, short, unsigned, char, etc).
- Generally, the operand will be an unsigned octal or an unsigned hexadecimal quantity, though this is not a firm requirement.

Ex: $07ff|_{16} = 0000\ 0111\ 1111\ 1111|_2$

$\sim(0000\ 0111\ 1111\ 1111) = 1111\ 1000\ 0000\ 0000$
 $= f800|_{16}$

$\sim(07ff) = f800$

Bases

- Normally , C assumes that integer constants are decimal (base 10)
- How can a computer tell whether a integer const is meant to be a **decimal, hexadecimal, or octal** value?
- A prefix 0x or 0X (Zero-exe) that you are specifying a hexadecimal value
- A prefix 0 (zero) means that you are writing in octal
Ex: 16 (decimal) = **020** (octal) = **0x10** (hexadecimal)
- In the binary code used internally by computers.
- To display an integer in octal use %o in hexadecimal %x.
- If you want to display prefixes **%#o , %#x , %#X**

Bases

```
#include<stdio.h>
main(){
    unsigned i= 0x5b3c;
        printf("Hex i = %x ~i = %x \n", i, ~i);
        printf("Dec i = %u ~i = %u \n", i, ~i);
} /* end of main */
```

```
#include<stdio.h>
main(){
    int x = 100;
        printf("hex = %x; oct = %o; dec = %d\n", x, x, x);
        printf("hex = %#x; oct = %#o; dec = %d\n", x, x, x);
}
hex = 64;   oct = 144;  dec = 100
hex =0x64; oct= 0114; dec= 100
```

Bit Level Operations

- Operations on bits and string of bits.

& : Bitwise AND

| : Bitwise OR

^ : Bitwise exclusive OR

~ : One's complement

<< : Left shift

>> : Right shift

- This a special feature of C programming language

a	0	0	1	1
b	0	1	0	1
a&b	0	0	0	1
a b	0	1	1	1
a^b	0	1	1	0
~a	1	1	0	0

String p	p >> 3	Action
01011010	00001011	P shifted right by 3 position. Left most bits filled with 0s
String q	q << 4	Action
11000110	01100000	Q shifted left by 4 positions. Right most bits filled with 0s

Bit Level Operations

The AND Gate

$$0 \& 0 = 0$$

$$1 \& 0 = 0$$

$$0 \& 1 = 0$$

$$1 \& 1 = 1$$

$$1010 \& 1100 \\ = 1000$$

Exclusive-OR, or XOR

$$0 \wedge 0 = 0$$

$$1 \wedge 0 = 1$$

$$0 \wedge 1 = 1$$

$$1 \wedge 1 = 0$$

$$1010 \wedge 1100 \\ = 0110$$

The OR Gate

$$0 | 0 = 0$$

$$1 | 0 = 1$$

$$0 | 1 = 1$$

$$1 | 1 = 1$$

$$1010 | 1100 \\ = 1110$$

The NOT Gate, or Inverter

$$\sim 0 = 1$$

$$\sim 1 = 0$$

$$\sim 10101 = 01010$$

Bit Level Operations

Adding Binary Numbers

$$0+0=0$$

$$1+0=1$$

$$1+1=10$$

$$\begin{array}{r} 11 \text{ (carry)} \\ 1010 \\ +1111 \\ \hline 11001 \end{array}$$

Binary Multiplication

$$1*1=1$$

$$1*0=0$$


$$0*1=0$$

$$\begin{array}{r} 101 \\ * 11 \\ \hline 101 \\ 1010 \\ \hline 1111 \end{array}$$

Shifting

a = 0110 1101 1011 0111 = 0x6db7

a<<6 = 0110 1101 1100 0000 = 0x6dc0



```
main(){
```

```
    unsigned a = 0xf05a;
```

```
    int b = a;
```

```
    printf(“%u %d\n”, a, b); 1111 0000 0101 1010 = 0xf05a
```

```
    printf(“%x\n”, a>>6); 0000 0011 1100 0001 = 0x3c1
```

```
    printf(“%x\n”, b>>6); 1111 1111 1100 0001 = 0xffc1
```

Masking

- Masking is a process in which a given bit pattern is transformed into another bit pattern by means of a logical bitwise operation.

Example:

Unsigned int a=0x6db7, b;

Extract the rightmost 6 bits and assign to b

b = a & 0x3f

a = 0110 1101 1011 0111

mask = 0000 0000 0011 1111

b = 0000 0000 0011 0111 = 0x37

The mask prevents the leftmost 10 bits from being copied from a to b

Example:

Initially $a=0x6db7$ assume 16-bit word size

- $a\&=0x7f \Rightarrow a=a \& 0x7f$ $a=$ $0x37$
- $a^{\wedge}=0x7f$ $0x6dc8$
- $a|=0x7f$ $0x6dff$
- $a\ll=5$ $0xb6e0$

Pointer to pointer

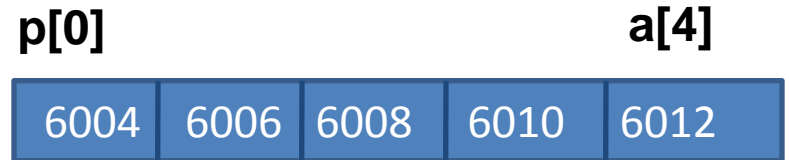
```
main(){
    static int a[] = {0,1,2,3,4};
    static int *p[] = {a, a+1, a+2, a+3, a+4};
    int **ptr;
    ptr=p;

    **ptr++;

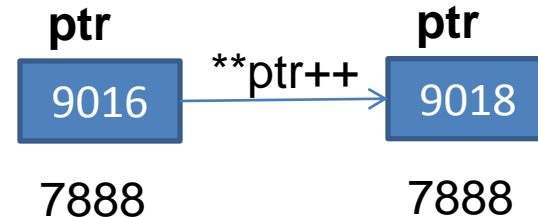
}
```



6004 6006 6008 6010 6012



9016 9018 9020 9022 9024



Recursive function

```
function fact(N){  
    if(N==1) return 1;  
    else return N*fact(N-1);  
}
```

- Let $N = 3$, the $3! = 3 \times 2 \times 1 = 6$.
- First time the function is called, $N = 3$ the value of N , but not of $\text{fact}(N-1)$, pushes N (value=3) on the stack, and calls itself for the second time with the value 2. ...
- Third time function calls itself with the value 1
- as $N=1$, the function returns 1. Since the value of $\text{fact}(1)$ is now known, it reverts back to its second execution by popping the last value (2) from the stack and multiplying it by 1.
- This operation gives the value of $\text{fact}(2)$, so the function reverts to its first execution by popping the next value (3) from the stack, and multiplying it with $\text{fact}(2)$, giving the value 6, that's the value the function finally returns.